

§ 1 Notation

The letters 'a', 'b', and 'c' with or without any premises will be used as the ^{arbitrary} names of simples. Elementary propositions will be written as a capital letter followed by any number of names of simple, separated by commas and enclosed by parenthesis, eg. $R(a, b)$.

The following things will be assumed:

1 all elementary propositions are propositions.

2 of (\bar{p}) is any group of propositions ' $N(\bar{p})$ ' is a proposition which is true if all the propositions in the group (\bar{p}) are false and ' $N(\bar{p})$ ' is false otherwise.

The letters 'p', 'q', and 'r' will be used as unspecified propositions. I now define

- D1 *1.01 $\lceil p \mid q \rceil \stackrel{df}{=} \lceil N(p, q) \rceil$
- D2 *1.02 $\lceil \sim p \rceil \stackrel{df}{=} \lceil p \mid p \rceil$
- D3 *1.03 $\lceil p \cdot q \rceil \stackrel{df}{=} \lceil \sim p \mid \sim q \rceil$
- D4 *1.04 $\lceil p \vee q \rceil \stackrel{df}{=} \lceil \sim (p \mid q) \rceil$
- D5 *1.05 $\lceil p \supset q \rceil \stackrel{df}{=} \lceil \sim p \vee q \rceil$
- D6 *1.06 $\lceil p \equiv q \rceil \stackrel{df}{=} \lceil (p \supset q) \cdot (q \supset p) \rceil$

I am using quasi-notation
as defined by Quine in
Mathematical Logic.

D7 Now let (\bar{p}) be all the propositions
which can be obtained by putting
in an appropriate name of some
simple in for the name of a
simple in a proposition. I define

$$\lceil (\exists \alpha) \phi \rceil \stackrel{df}{=} \lceil N(N(\bar{p})) \rceil$$

On this, ϕ is the name of an arbitrary

sign with is the result of putting
 α for the name of a simple in
 a proposition.

$$D8 * 1.08 \quad \lceil \alpha \rceil \stackrel{df}{=} \lceil \sim (\exists \alpha) \sim \rceil$$

D9 Let (\bar{p}) be all the propositions which
 can be obtained by putting in
 the names of any two appropriate
 simples in the place of two names
 of simples in a proposition.

$$1.09 \quad \lceil (\exists \beta)(\exists \alpha) \phi \rceil \stackrel{df}{=} \lceil N(N(\bar{p})) \rceil, \text{ etc}$$

From now on, the Greek letters
 $\phi, \chi,$ and λ will be used
 for expressions which will
 become propositions if a number
 of variables ($\alpha, \beta, \gamma, \dots$) are
 replaced by the names of
 appropriate simples. The occurrence
 of any number of sentential
 quantifiers is defined by D7 and
 D9. The occurrence of universal
 quantifiers is defined by D8.

The closure of ϕ is the proposition
which results from applying
in alphabetic order $(x, y, z, x', y', z', \dots)$
the corresponding universal
quantifier for every variable
in the expression. That the
closure of ϕ is a theorem will
be written $\vdash \phi$.

§2 Primitive Propositions.

I have so far only made definitions. Before anything can be proved, something must be assumed. The following are descriptions of all the propositions I will assume to be logically true. They come from Mathematical Logic by Quine. I have modified the notation some, and the generation of ^{all} propositions from elementary propositions has been made explicit.

*100 ~~WVA~~ If ϕ is tautologous, by a truth table test $\vdash \phi$.

*101 ~~WVA~~ $\vdash \lceil (\alpha)(\phi \supset \psi) \supset (\alpha)\phi \supset (\alpha)\psi \rceil$

*102 ~~WVA~~ If ϕ' is like ϕ except for certain free occurrences of α' wherever ϕ contains free occurrences of α , then $\vdash \lceil (\alpha)\phi \supset \phi' \rceil$

*104 ~~iff~~ iff $\vdash \lceil p \supset q \rceil$ and $\vdash p$, then $\vdash q$.

~~\rightarrow 5 $\vdash \lceil N(p) \rceil$ and~~

~~\rightarrow 4 iff $\vdash \lceil p \supset N(p) \rceil$ and $\vdash p$, and
for any (p) , then $\vdash \lceil \neg q \rceil$~~

*105 ~~iff~~ $\vdash \lceil N(\bar{p}) \supset \sim p \rceil$ where
 p is in (\bar{p}) .

*106 iff for any p in (\bar{p}) , $\vdash \sim p$, then
~~iff~~
 $\vdash \lceil N(\bar{p}) \rceil$

*107 iff for any p in (\bar{p}) there
is a q in (\bar{q}) such that $\vdash \lceil q \supset p \rceil$
and $\vdash \lceil N(\bar{p}) \supset N(\bar{q}) \rceil$ (.)

*108 $\vdash \lceil N(\bar{p}) \cdot N(\bar{q}) \supset N(\bar{F}) \rceil$
where any r in (\bar{r}) is also in (\bar{p}) or (\bar{q}) .

§ 7 Identity & Quantification of functions

Let (\bar{p}) be the class of all elementary propositions of the form $\neg(\phi \equiv \psi)$ where ϕ is like ψ except for containing a where ψ contains b , then

D10 $\neg(a=b) \stackrel{df}{=} \neg N(\bar{p})$

Let (\bar{p}) be the class of all propositions ψ which contain a proposition in a number of places (30) such that ~~the same argument the~~

D11 $\neg(\exists \phi) \psi \stackrel{df}{=} N(N(\bar{p}))$

D12 $\neg(\exists \phi) \neg \psi \stackrel{df}{=} N(\exists \phi) \neg \psi$

*180 If ϕ and ψ are the same except that ϕ contains a where ψ contains b , then

$$\vdash \lceil a=b \supset \phi \equiv \psi \rceil$$

Proof: $\vdash \lceil a=b \supset N(\bar{p}) \rceil$ *100 + D 10
 $\vdash \lceil N(\bar{p}) \supset \sim(\sim(\phi \equiv \psi)) \rceil$ D10 + *105
 $\vdash \lceil \sim(\sim(\phi \equiv \psi)) \supset \phi \equiv \psi \rceil$ *100
 $\vdash \lceil a=b \supset \phi \equiv \psi \rceil$ *112

~~*181 If ϕ and ψ are alike except that ϕ contains a where ψ contains b , then~~

~~$$\vdash \lceil (\beta)(\beta = a \supset \psi) \equiv \phi \rceil$$~~

~~$$\vdash$$~~

*181 $\vdash \lceil a=a \rceil$

Let (p) be the set of all elements which contain a . ~~of p~~
~~for any $p \in (p)$, $\sim(p \equiv p)$~~
 is contradictory, so

$\vdash \lceil N(\bar{p}) \rceil$ *106
 $\vdash \lceil N(\bar{p}) \supset a=a \rceil$ D 10
 $\vdash \lceil a=a \rceil$ *112

$$*182 \vdash \lceil a=b \supset b=a \rceil$$

$$\vdash \lceil a=b, \supset N(\bar{p}) \rceil \quad D10$$

Let (\bar{q}) be the set of all propositions of the form $\lceil \sim (\psi \equiv \phi) \rceil$.
 $\emptyset \neq \bar{q}$ as in D10. Then for any $p \in (\bar{p})$ there is a q in (\bar{q}) such that $q \supset p$.

$$\vdash \lceil N(\bar{p}) \supset N(\bar{q}) \rceil \quad *107$$

$$\vdash \lceil N(\bar{q}) \supset .b=a \rceil \quad D10 *100$$

$$\vdash \lceil a=b, \supset .b=a \rceil \quad *112$$

$$*183 \vdash \lceil b=a \supset a=b \rceil$$

$$*182$$

$$*184 \vdash \lceil a=b \equiv b=a \rceil$$

~~$\vdash \lceil a=b \supset b=a \rceil$~~
 ~~$\vdash \lceil b=a \supset a=b \rceil$~~
 ~~$\vdash \lceil a=b \equiv b=a \rceil$~~

$$\vdash \lceil [182] \supset . [183] \supset [182], [183] \rceil \quad *100$$

$$\vdash \lceil [182] \rceil \quad *182$$

$$\vdash \lceil [183] \supset [182], [183] \rceil \quad *104$$

$$\vdash \lceil 183 \rceil \quad *183$$

$$\vdash \lceil [182], [183] \rceil \quad *104$$

$$\begin{aligned} &\vdash \lceil [182], [183] \supset [184] \rceil \quad D6 + *100, \\ &\vdash \lceil [184] \rceil \quad *104 \end{aligned}$$

$$*185 \quad \vdash \lceil a=b, b=c, \supset, a=c \rceil$$

$$\vdash \lceil a=b, b=c, \supset, N(\bar{p}), N(\bar{q}) \rceil \quad *100 + D6,$$

$$\vdash \lceil N(\bar{p}), N(\bar{q}), \supset N(\bar{r}) \rceil$$

where (\bar{r}) is the contr. of a

simple proposition . . .

$$\vdash \lceil N(\bar{r}) \supset, a=c \rceil \quad *100 + D6,$$

$$\vdash \lceil a=b, b=c, \supset, a=c \rceil \quad *112,$$

$$*186 \quad \vdash \lceil (\lambda) (\lambda = \lambda) \rceil$$

let (\bar{p}) be all propositions of the form $\lceil \sim(a=a) \rceil$

whenever (\bar{p}) in (\bar{p}) is $\sim p$ is a theorem (*151)

$$\vdash \lceil N(\bar{p}) \rceil \quad *106$$

$$\vdash \lceil N(\bar{p}) \supset \sim(\exists x) \sim(x=x) \rceil \quad *100 + D7$$

$$\vdash \lceil \sim(\exists x) \sim(x=x) \supset (\lambda) (x=x) \rceil \quad *100 + D8$$

$$\vdash \lceil N(\bar{p}) \supset (\lambda) (x=x) \rceil \quad *112$$

$$\vdash \lceil (\lambda) (x=x) \rceil \quad *104$$

$$D1 \quad \lceil \phi \mid \psi \rceil \stackrel{df}{=} \lceil N(\phi, \psi) \rceil$$

$$D2 \quad \lceil \sim \phi \rceil \stackrel{df}{=} \lceil N(\phi) \rceil$$

$$D3 \quad \lceil \phi, \psi \rceil \stackrel{df}{=} \lceil \sim \phi \mid \sim \psi \rceil$$

$$D4 \quad \lceil \phi \vee \psi \rceil \stackrel{df}{=} \lceil \sim(\phi \mid \psi) \rceil$$

$$D5 \quad \lceil \phi \supset \psi \rceil \stackrel{df}{=} \lceil \sim \phi \vee \psi \rceil$$

$$D6 \quad \lceil \phi \equiv \psi \rceil \stackrel{df}{=} \lceil (\phi \supset \psi) \cdot (\psi \supset \phi) \rceil$$

D7 $\lceil (\exists x)\phi \rceil \stackrel{df}{=} \lceil \sim(N(\bar{p})) \rceil$, where (\bar{p}) is all the propositions which can be obtained by putting in the name of some appropriate simple in for the name of a simple in a proposition, and ϕ is that proposition with x substituted for the name of that simple.

$$D8 \quad \lceil \sim(x) \rceil \stackrel{df}{=} \lceil \sim(\exists x)\sim \rceil$$

*100 If ϕ is tautologous, $\vdash \phi$.

*105 If for any ϕ in (\bar{p}) there is a ψ in (\bar{q}) such that $\vdash \psi \supset \phi$, then
 $\vdash \lceil N(\bar{p}) \supset N(\bar{q}) \rceil$

*106 $\vdash \lceil (\alpha)(\phi \supset \psi) \supset (\alpha)\phi \supset (\alpha)\psi \rceil$

Proof

$\vdash \lceil (\alpha)(\phi \supset \psi) \supset \sim(\exists \alpha) \sim(\phi \supset \psi) \rceil$ *100 + D8

$\vdash \lceil \sim(\exists \alpha) \sim(\phi \supset \psi) \supset N(\bar{p}) \rceil$ *100 + D7

where (\bar{p}) is all the propositions
obtained from $\sim(\phi \supset \psi)$ by
substituting the name of an appropriate
variable for \sim

$\vdash \lceil N(\bar{p}) \supset N(\bar{q}) \rceil$

*105

where